# CSC3100 Data Structures
# Tutorial 3

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https://github.com/I-am-Future

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- Complexity for recursion and divide-and-conquer (Prof. Fang’s lec 6, Prof. Yu’s lec 5)
  - Concepts
  - Practice Problems
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# 1. Asymptotic Analysis-Concepts

- Evaluate the "efficiency" of an algorithm.

- Commonly used notations:
  - Big-Oh notation: measure the upper bound complexity.
  - Big-Omega notation: measure the lower bound complexity.
  - Big-Theta notation: just there! Upper & lower bounds meet!

# Upper & lower bound? (1)

- For some algorithm f(n), assume we can prove the greens:
- $O(1)$  $O(\log n)$  $O(n)$  $O(n \cdot \log n)$  $O(n^2)$  $O(n^3)$  $O(2^n)$

^^^^^^^^^

(minimum upper bound)

---

- $\Omega(1)$  $\Omega(\log n)$  $\Omega(n)$  $\Omega(n \cdot \log n)$  $\Omega(n^2)$  $\Omega(n^3)$  $\Omega(2^n)$

^^^^^^^^^

(maximum lower bound)

- The minimum upper bound and maximum lower bound meet
- => f(n) is $\Theta(n \cdot \log n)$

# Upper & lower bound? (2)

- For some algorithm f(n), assume we can prove the greens:
- $O(1)$ $O(\log n)$ $O(n)$ $O(n \cdot \log n)$ $O(n^2)$ $O(n^3)$ $O(2^n)$

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^^^^^^^^^

(maximum lower bound)

- The minimum upper bound and maximum lower bound don't meet
- => we cannot conclude f(n) is $\Theta(n \cdot \log n)$. Usually, we only say f(n) is $O(n^2)$

# Upper & lower bound? (3)

- For some algorithm f(n), assume we can prove the greens:
- $O(1)$ $O(\log n)$ $O(n)$ $O(n \cdot \log n)$ $O(n^2)$ $O(n^3)$ $O(2^n)$

^^^^

(minimum upper bound)

---

- $\Omega(1)$ $\Omega(\log n)$ $\Omega(n)$ $\Omega(n \cdot \log n)$ $\Omega(n^2)$ $\Omega(n^3)$ $\Omega(2^n)$

^^^^^^^^^

(maximum lower bound)

- Upper bound is smaller than lower bound: such case NEVER happen!!!

# Rules: help you calculate complex one

For Big-Oh and Big-Omega, all of them are valid:

1. Polynomial Rule: Only the biggest matter
2. Product Rule: the big multiplies the big
3. Sum Rule: the bigger of the two big
4. (Log Rule): Log only beats constant
5. (Exponential Rule): Exponential beats power functions

# 1. Asymptotic Analysis-Practice Problems

```
for (int i = 0; i < n; i++) {
    // Some O(1) operation
}
```

---

```
for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
        // Some O(1) operation
    }
}
```

# 1. Asymptotic Analysis-Practice Problems

```
for (int i = 0; i < n; i++) {
    // Some O(1) operation
}
```

$O(n)$

```
for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
        // Some O(1) operation
    }
}
```

$O(n^2)$

# 1. Asymptotic Analysis-Practice Problems

```
for (int i = 0; i < n; i+=(n/2)) {
    // Some O(1) operation
}
```

---

```
while (n > 0) {
    if (n % 2 == 1)
        res = res * a;
    a = a * a;
    n = n / 2;
}
```

# 1. Asymptotic Analysis-Practice Problems

```
for (int i = 0; i < n; i+=(n/2)) {
    // Some O(1) operation
}
```

$O(1)$

---

```
while (n > 0) {
    if (n % 2 == 1)
        res = res * a;
    a = a * a;
    n = n / 2;
}
```

This algorithm is called “Quick Power”

$O(\log n)$

**More Info:** https://www.rookieslab.com/posts/fast-power-algorithm-exponentiation-by-squaring-cpp-python-implementation

# 1. Asymptotic Analysis-Practice Problems

check and prove $g(n) = (0.1n^2 + n\log n) \cdot (n\log n + \sqrt{n}) = \Theta(n^3 \cdot \log n)$.

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check and prove $g(n) = (0.1n^2 + n\log n) \cdot (n\log n + \sqrt{n}) = \Theta(n^3 \cdot \log n)$.

**Idea：**

1. "Expand" the g(n) to 4 terms
2. Prove $g(n) = O(n^3 \cdot \log n)$ : Apply rule 3: Sum rule, only choose biggest among all.
3. Prove $g(n) = \Omega(n^3 \cdot \log n)$ : Apply rule 3: Sum rule, only choose biggest among all.
4. Done!

**Two key points:**

1. Apply the rule to simplify the problem;
2. When proving $\Theta(\cdot)$, we need to prove $O(\cdot)$ and $\Omega(\cdot)$ together.

# 2. Complexity for recursion and divide-and-conquer - Concepts

- To calculate the complexity for recursion and divide-and-conquer algorithm:

- Step 1: Get the recursive expression (looks like: $g(n) = g(n-1) + O(n)$, $g(n) = g(n/2) + O(n)$ )
- Step 2:
  - Method 1: Unfold $g(n)$ to $g(1)$ by hand and get the answer.
  - Method 2: Master theorem--

- Recurrence: $T(n) \leq a \cdot T(n/b) + O(n^d)$
- An algorithm that divides a problem of size n into a subproblems, each of size n / b

$$T(n) = \begin{cases} O(n^d \log n) & \text{if } a = b^d \\ O(n^d) & \text{if } a < b^d \\ O(n^{\log_b(a)}) & \text{if } a > b^d \end{cases}$$

**a**: number of subproblems (branching factor)
**b**: factor by which input size shrinks (shrinking factor)
**d**: need to do $O(n^d)$ work to create subproblems + "merge" their solutions

From Prof. Fang's slides – lec 6 – page 19.

**4-1 *Recurrence examples***

Give asymptotic upper and lower bounds for $T(n)$ in each of the following recurrences. Assume that $T(n)$ is constant for $n \le 2$. Make your bounds as tight as possible, and justify your answers.

**a.** $T(n) = 2T(n/2) + n^4$.

**b.** $T(n) = T(7n/10) + n$.

**c.** $T(n) = 16T(n/4) + n^2$.

**d.** $T(n) = 7T(n/3) + n^2$.

**e.** $T(n) = 7T(n/2) + n^2$.

**f.** $T(n) = 2T(n/4) + \sqrt{n}$.

**g.** $T(n) = T(n-2) + n^2$.

**Let you be familiar with it!**

- Recurrence: $T(n) \le a \cdot T(n/b) + O(n^d)$
- An algorithm that divides a problem of size n into a subproblems, each of size n / b

$$T(n) = \begin{cases} O(n^d \log n) & \text{if } a = b^d \\ O(n^d) & \text{if } a < b^d \\ O(n^{\log_b(a)}) & \text{if } a > b^d \end{cases}$$

**a**: number of subproblems (branching factor)
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Problem 4-1 (page 107, 3rd edition)

# Question a-f: Use Master's Theorem.

**a.** By master theorem, $T(n) = \Theta(n^4)$.

**b.** By master theorem, $T(n) = \Theta(n)$.

**c.** By master theorem, $T(n) = \Theta(n^2 \lg n)$.

**d.** By master theorem, $T(n) = \Theta(n^2)$.

**e.** By master theorem, $T(n) = \Theta(n^{\lg 7})$.

**f.** By master theorem, $T(n) = \Theta(\sqrt{n} \lg n)$.

# Question g: Expand the recursion

- $T(n) = T(n-2) + n^2$
- $\quad = T(n-4) + (n-2)^2 + n^2$
- $\quad = T(1) + 3^2 + \ldots\ldots + (n-4)^2 + (n-2)^2 + n^2$ (If n is odd)
- $\quad = T(2) + 4^2 + \ldots\ldots + (n-4)^2 + (n-2)^2 + n^2$ (If n is even)

- By the sum of the squares formula, We know that T(n) is $\Theta(n^3)$.
- Some small tricks in our case:
  - Even n: use formula with n=2k in it.
  - Odd n: use sum difference (total sum-even sum)

$1^2+2^2+3^2+4^2+ \ldots +n^2$

Derivation:

$$\frac{n(n+1)(2n+1)}{6}$$

# Question g: Expand the recursion (details)

- Tricks to calculate:
- - Sum of even squares
- - Sum of odd squares

- $T(n) = 2^2 + 4^2 + \cdots \cdots + n^2$

let $n = 2k$:

$$T(n) = (2\times1)^2 + (2\times2)^2 + \cdots\cdots + (2k)^2$$
$$= 2^2(1^2 + 2^2 + \cdots\cdots + k^2)$$
$$= 4 \cdot \frac{k(k+1)(2k+1)}{6}$$
$$= \frac{n(n+2)(n+1)}{6} = \Theta(n^3)$$

- $T(n) = 1^2 + 3^2 + \cdots\cdots + n^2$

$$= \underbrace{(1^2 + 2^2 + \cdots + n^2)}_{\text{known}} - \underbrace{(2^2 + 4^2 + \cdots + (n-1)^2)}_{\text{known above!}}$$

# **Look back:** An example that upper/lower bounds does not meet (Will not be in exam)

## Upper & lower bound? (2)

- For some algorithm f(n), assume we can prove the greens:
- $O(1)$ $O(\log n)$ $O(n)$ $O(n \cdot \log n)$ $O(n^2)$ $O(n^3)$ $O(2^n)$

^^^^^

(minimum upper bound)

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- $\Omega(1)$ $\Omega(\log n)$ $\Omega(n)$ $\Omega(n \cdot \log n)$ $\Omega(n^2)$ $\Omega(n^3)$ $\Omega(2^n)$

^^^^^^^^^

(maximum lower bound)

- The minimum upper bound and maximum lower bound don't meet
- => we cannot conclude f(n) is $\Theta(n \cdot \log n)$. Usually, we only say f(n) is $O(n^2)$

For expression like:

$T(n) = T(7n/10) + \log(n)$

Cannot use Master Theorem.

--But we can do scaling (放缩)

--So we can find a good upper bound, O(n);

--And a good lower bound, Ω(log n)!

(Will not be tested. Just for fun!)

In Exercise 4-1 (b), we know:

$$T(n) = T\left(\frac{7n}{10}\right) + n \rightarrow O(n)$$

Question

But, if expression is

$$T(n) = T\left(\frac{7n}{10}\right) + \log(n)$$ ← We can't use Master Theorem

① $T(n) = T\left(\frac{7n}{10}\right) + \log(n) < T\left(\frac{7n}{10}\right) + n \rightarrow \underbrace{O(n)}$

upper bound of this

also upper bound of original

By Master Theorem

② $T(n) = T\left(\frac{7n}{10}\right) + \log(n) > T\left(\frac{7n}{10}\right) + n^0 \rightarrow \underbrace{\Omega(\log n)}$

lower bound of this

also lower bound of original

Conclusion:

By our knowledge so far, we can only conclude

$$T(n) = T\left(\frac{7n}{10}\right) + \log(n) = O(n), \Omega(\log n)$$

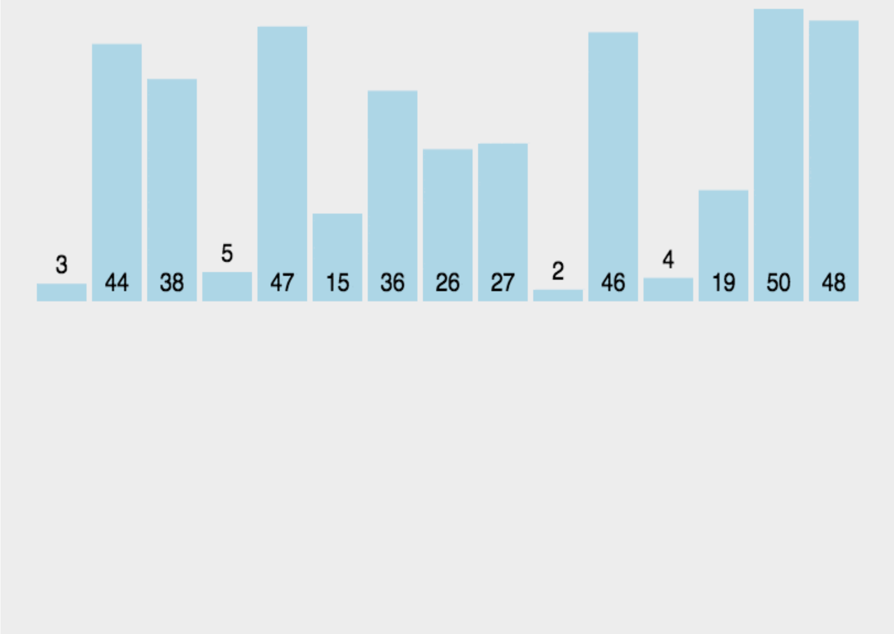
# Coding Question: Insertion Sort

- We provide three languages (C++/Python/Java) of code, with corresponding reference answer on Blackboard.

- The sorting function should **return void**, i.e., You need to modify in-place.
- The problem is not difficult, but you need to think about “edge cases”.

# How It Works:

1. **Start with the second element** (index 1) in the array, treating the first element as already sorted.
2. **Compare** the current element with the previous elements.
3. **Insert** the current element into the correct position by shifting the larger elements one position to the right.
4. **Repeat** the process for each of the elements in the array.




**Insertion sort** From Prof. Fang's slides – lec 4 – page 4.

- A simple algorithm for <u>a small number of elements</u>
- Similar to sort a hand of cards
  - Start with an empty left hand
  - Pick up one card and insert it into the correct position
  - To find the correct position, compare it with each of the cards in the hand, from right to left
  - The cards in the left hand are sorted

# 3. Q & A